

# A set theoretic framework for enumerating matches in surveys and its application to reducing inaccuracies in vehicle roadside surveys

Richard G. Clegg

*Department of Mathematics, University of York,  
York, YO10 5DD, United Kingdom*

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## Abstract

This paper describes a framework for analysing matches in multiple data sets. The framework described is quite general and can be applied to a variety of problems where matches are to be found in data surveyed at a number of locations (or at a single location over a number of days). As an example, the framework is applied to the problem of false matches in licence plate survey data. The specific problem addressed is that of estimating how many vehicles were genuinely sighted at every one of a number of survey points when there is a possibility of accidentally confusing two vehicles due to the nature of the survey undertaken.

In this paper, a method for representing the possible *types of match* is outlined using set theory. The phrase *types of match* will be defined and formalised in this paper. A method for enumerating  $\mathcal{M}_n$ , the set of all types of match over  $n$  survey sites, is described. The method is applied to the problem of correcting survey data for false matches using a simple probabilistic method. An algorithm is developed for correcting false matches over multiple survey sites and its use is demonstrated with simulation results.

*Key words:* Traffic, Transportation, Applied Probability, Uncertainty Modelling

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## 1. Introduction

In the analysis of roadside survey data, it is often desirable to analyse matches between several data sets simultaneously. For example, we might wish to answer questions of the general type “How many drivers are seen at point A, point B and point C?” or “How many vehicles are seen on all five survey days?” This paper attempts to create a general framework for the analysis of matching between data from more than two surveys. The framework

is then applied to the specific case of false matching in partial licence plate surveys (that is non-matches which are mistaken for matches because only part of the licence plate is observed). It should be stressed throughout that the framework outlined is applicable to any data series where matches are sought between two or more distinct data sets. While the work is placed in the context of licence plate surveys (and further in the context of licence plate surveys using a specific type of British licence plate) the results are much more general than this.

Licence plate surveys are commonly used in the study of traffic systems, particularly when measurements of the same vehicle are required more

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*Email address:* richard@richardclegg.org (Richard G. Clegg).

than one point (for example, calculating travel times or the routes of vehicles). Although automated techniques are becoming more common (GPS, toll-tags and automated recognition cameras) the manual licence plate survey remains an important tool for the road transport engineer. If a road with a high volume of traffic is being surveyed then it is often the case that only part of the licence plate is recorded. When this is the case, the possibility of spurious matches occurs. To take an example, standard British licence plates used to be of the following form: single letter, three digits, three letters: e.g. A123BCD. This form will be used throughout the paper, however, it must be stressed that this method would work with partial observations of any type given the assumptions stated later. If a surveyor only recorded the first letter and three digits, then a vehicle A123ABC would not be distinguished from a vehicle A123XYZ since the disambiguating information (the final three letters) would not be recorded.

While the chances of such a false match are low, quite often the combinatorics of the problem means that the actual recorded number of false matches remains high. To mathematicians, this is familiar as the celebrated *Birthday Paradox*. The Birthday Paradox asks the question “How many people must we have in a room before we might expect that two share the same birthday?” Intuitively, we might expect this to be quite a high number (since it is unlikely that any two people share a birthday). However, the number of pairs of people in a room goes up with the square of the number of people in the room  $(n^2 - n)/2$ . If we made the assumption that the chance of two randomly selected people sharing a birthday is one in 365 then we only need twenty three people in the room before it becomes likely (probability above 50%) that two will share a birthday. Combinations in multiple point surveys work similarly. If we had two survey sites, each with one thousand observations then this is one million pairs of observations. If the chances of a false match in a given pair are only one in a ten thousand, we will still get (on average) one hundred false matches. This could well be larger than the actual number of genuine matches in the data set and will certainly be a significant bias.

This paper attempts to provide a sound theo-

retic backing (using the well-known framework of set theory) to matching problems across multiple data sites. In section two, a general background of matching problems in licence plate data is given to put the problem into context within the transport field. In section three, the concept of *types of match* is formalised using the standard set theoretic concept of an equivalence class. In section four, a simple method is given for constructing the set  $\mathcal{M}_n$ , the set of every possible type of match across  $n$  survey sites. In section five, partial ordering is introduced to apply the problem to false matches due to incomplete observations. In section six, an algorithm is given for correcting false matches using the framework developed in sections three, four and five. Finally, in section seven, computational results are given on artificially generated survey data. The work in this paper can be found in a much expanded form in (1, Chapter Four) and an example of the method being used on real road traffic data is found in (1, Chapter Five). The set theory used in this paper is extremely simple (just the concepts of equivalence class and partial order are necessary) and would be covered in any standard text on the subject, for example (2).

## 2. The false match problem in licence plate data at multiple sites

Throughout this paper, the examples are given using an old form of British licence plate — it should be stressed that this is not necessary for this framework and is done purely for the sake of example. The work described here assumes nothing about the nature of the individuals being observed other than the restrictions described in Definition 20. Similarly, when the phrase observation sites is used throughout this paper, this can mean either geographically distinct observation sites or a single geographical location observed for a number of days or any combination of times and locations. (In the work which motivated this research, the experimenters were interested in finding vehicles which travelled between three distinct geographical locations on two consecutive days. This would count as six observation sites in the terminology

used here.) Note that no time information is used here although time information is often available for such surveys. It is hoped that a future improvement to this method will make use of time information about observations to reduce uncertainties.

It is often the case that on-street traffic surveys collect partial vehicle licence plate information. [The reason for collecting partial rather than full licence plate information is that the recording and transcription of the data is often done manually and time constraints would preclude recording a full plate.] This information can then be used to reconstruct travel times and to infer route information about drivers. In partial plate data, however, problems can occur from *false matches* as discussed above. Of course, false matches could also occur through recording or transcription errors. While this paper will not discuss these problems, it is in principle possible to extend this framework to cover recording and transcription errors.

In the case of two survey sites and no recording or transcription errors the situation is relatively clear. If our data shows that a match occurs between two observations (one from each site) then, this must mean that either the same vehicle has been observed at both, or that two different vehicles have been observed which happened to have the same partial licence plate. At multiple sites the situation is much more complex. An apparent match at four survey points may be any of the following: a true match (the same vehicle seen at all four points); a different vehicle at each of the four points which (by coincidence) have the same partial plate; a vehicle at survey point one and two which has the same partial plate as a second vehicle at survey points three and four; or any other of fifteen total possibilities. The problem becomes more difficult as the number of sites increases. Indeed it is not immediately clear how to enumerate the number of ways in which a match as described above can occur over multiple data sites. This issue is not a trivial one. In real licence surveys, the number of false matches is often greater than the number of true matches. In (1, Chapter 5) two survey sites with a flow of approximately one thousand vehicles at each were found to have ninety observed matches between vehicles despite the fact that (given the positioning of the sites) it

would be extremely unlikely for any drivers at all to travel between them.

A number of researchers have approached the false matching problem for licence plates. An early approach for two sites is given by (3) which uses a simple probabilistic correction. Several methods are described in (4) including the possibility of two point matches between vehicles observed at pairs of sites selected from several survey sites (for example entering and leaving a cross-roads). A graphical procedure for visualising matches based upon journey time between two sites is given by (5). Methods in this paper are useful for any analysis of data in which time between observations is a factor. Further refinements for site pairs, including a maximum likelihood method based upon assumptions about travel time distribution are given in (6) and (7). However, all of these methods concentrate on matches between pairs of sites and the majority of them also assume that journey time information can be used to aid in finding false matches, which is not the case if, for example, we are interested in correcting false matches at the same site over different days. The method described in this paper concentrates on matches between observations at more than two sites, particularly where journey time information is not available or cannot be used.

It should be emphasised again that, while this work is presented within the context of licence plate surveys (indeed within the context of licence plate surveys on a specific type of British licence plate) the results presented are extremely general. These results would be applicable to any type of survey data where individuals are sought in more than two data sets and where a possibility of confusion between observations of individuals exists. Applications for this technique are being sought in other areas such as DNA matching and suggestions for suitable data sets would be welcomed by the author.

### **3. Equivalence classes for representing types of match**

In this section, notation is described, with examples, to describe a mathematical framework for

investigating matches in multiple data sets. For the convenience of the reader the notation used throughout this paper is gathered here for reference and defined as it occurs throughout the paper. In general bold lower case  $\mathbf{x}$  is used to indicate a tuple (ordered set). Upper case  $M$  is used to indicate a set and bold upper case  $\mathbf{M}$  is used to indicate a set of sets. Caligraphic lettering  $\mathcal{S}$  is used to indicate higher order entities such as sets of tuples or sets of sets of sets.

The following specific notation is used.

- $n$  — the number of sites under investigation.
- $\#M$  — the number of members of set  $M$ .
- $S_i$  — the set of observations at site  $i$ . See Definition 1.
- $\mathbf{y}$  — a tuple of observations, one from each site. See Definition 2
- $\mathcal{S}$  — the set of all possible tuples of observations. See Definition 3
- $\mathbf{M} = \{M_1, M_2, \dots, M_m\}$  — a *type of match*. See Definition 5.
- $\mathcal{M}_n$  — the set of all types of match for  $n$  sites. See Definition 6.
- $C(\mathbf{y})$  — the *type of match* of a tuple of observations  $\mathbf{y}$ . See Definition 7.
- $\mathbf{A}_n$  — the set of sets  $\{\{1, 2, \dots, n\}\}$  representing the same observation across all sites. See Definition 9.
- $\mathbf{y}^*$  — the tuple of *partial observations* from the tuple  $\mathbf{y}$ . See Definition 11.
- $\mathcal{S}^*$  — the set of all such *partial observations*. See Definition 11.
- $x(\mathbf{y}, \mathbf{M})$  — the *exact matching function* for the tuple  $\mathbf{y}$ . See Definition 13.
- $X(\mathbf{M})$  — the exact matching count for the set  $\mathcal{S}$ . See Definition 14.
- $r(\mathbf{y}, \mathbf{M})$  — the *relaxed matching function* for the tuple  $\mathbf{y}$ . See Definition 15.
- $R(\mathbf{M})$  — the relaxed matching count for the set  $\mathcal{S}$ . See Definition 16.
- $T(M)$  — the number of observations which are the same across all sites in the set  $M$ . See Definition 18.
- $p(i)$  — the probability that  $i$  distinct individuals, different in a full observation, are the same in a partial observation. See Definition 20.

**Definition 1** *Let  $n$  be the number of observation sites and let  $S_i$  be the set of observations at the  $i$ th*

*such site.*

Consider the following toy example with three sites ( $n = 3$ ),

$$\begin{aligned} S_1 &= \{\mathbf{A123XYZ}, \mathbf{C789ABC}\} \\ S_2 &= \{\mathbf{A123XYZ}, \mathbf{A123XDR}, \mathbf{D555SDD}\} \\ S_3 &= \{\mathbf{C789ABC}, \mathbf{A123XYZ}\}. \end{aligned}$$

In passing, it should be noted that a formal requirement for something to be a set is that its members are distinct. If this formal requirement is not met then each member of the set could be tagged by a unique number which is not considered in later equality relations. This is a technicality which will not be mentioned again and does not affect what follows.

**Definition 2** *A tuple of observations  $\mathbf{y} = (y_1, \dots, y_n)$  is an  $n$ -tuple consisting of one member of each set of observations — that is,  $y_i \in S_i$  for all  $i$ .*

Continuing the previous example,

$$\mathbf{y} = (\mathbf{A123XYZ}, \mathbf{A123XYZ}, \mathbf{C789ABC})$$

is the tuple formed by taking the first observation from each set.

**Definition 3** *The set of all tuples of observations  $\mathcal{S}$  in the data is the set of all such  $\mathbf{y}$  which can be formed from the sets  $S_1, \dots, S_n$ . This is clearly the cartesian product given by*

$$\mathcal{S} = S_1 \times S_2 \cdots \times S_n.$$

So, in the example framework given before, then the set  $\mathcal{S}$  has twelve members and is given by

$$\begin{aligned} \mathcal{S} = \{ &(\mathbf{A123XYZ}, \mathbf{A123XYZ}, \mathbf{C789ABC}), \\ &(\mathbf{A123XYZ}, \mathbf{A123XYZ}, \mathbf{A123XYZ}), \\ &\dots (\mathbf{C789ABC}, \mathbf{D555SDD}, \mathbf{A123XYZ}) \}. \end{aligned}$$

Considering, the members of  $\mathcal{S}$  it is obvious that  $(\mathbf{A123XYZ}, \mathbf{A123XYZ}, \mathbf{A123XYZ})$  is the type of observation which is most of interest, the same individual observed across all three sites under investigation. Also, in some way, the tuples  $(\mathbf{A123XYZ}, \mathbf{A123XDR}, \mathbf{A123XYZ})$  and  $(\mathbf{C789ABC}, \mathbf{A123XDR}, \mathbf{C789ABC})$  are in some way structurally similar (they match at sites one and three) and both are structurally different to

(A123XYZ, A123XYZ, C789ABC).

This structural similarity will now be formalised by using the concept of a *type of match*.

**Definition 4** Two  $n$ -tuples of observations  $\mathbf{y} = (y_1, \dots, y_n)$  and  $\mathbf{z} = (z_1, \dots, z_n)$  are the same type of match ( $\mathbf{y} \sim \mathbf{z}$ ) if they are equal at the same sites (and different at the same sites). Formally,

$$\mathbf{y} \sim \mathbf{z} \text{ if and only if } (y_i = y_j) \Leftrightarrow (z_i = z_j) \\ \text{for all } i, j \in \{1, 2, \dots, n\}.$$

Note that, for simplicity the limits  $i, j \in \{1, 2, \dots, n\}$  on indices will usually be omitted where, as in this case, they are obvious. It can trivially be shown that the relation defined by  $\sim$  meets the requirements of an equivalence relation in set theory.

#### 4. The set of every type of match

Having formalised the concept of when two sets of observations are the same type of match, the next step is to introduce an entity which can represent the type of match of a given tuple of observations. This is simply achieved using partitions of the first  $n$  integers. A partition of the first  $n$  integers is a set of sets  $\mathbf{M} = \{M_1, M_2, \dots, M_m\}$  such that each integer from one to  $n$  is in one and only one of the sets  $M_1 \dots M_m$ . (In the literature, these  $M_i$  are often referred to as blocks.) Any  $n$ -tuple of observations is related to some such  $\mathbf{M}$  by the relation given in Definition 7.

**Definition 5** A type of match is a partition  $\mathbf{M}$  of the first  $n$  integers which is used to represent the structure of matches within an  $n$ -tuple of observations  $\mathbf{y}$ . The relationship between  $\mathbf{M}$  and  $\mathbf{y}$  is given by Definition 7.

Considering the first three integers, then  $\{\{1, 2, 3\}\}$ ,  $\{\{1, 2\}, \{3\}\}$  and  $\{\{1\}, \{2\}, \{3\}\}$  are among the possible partitions.

**Definition 6** The set  $\mathcal{M}_n$  is the set of all possible partitions of the first  $n$  integers. This can be used to represent any possible type of match over  $n$  observation sites.

For one site only the partition  $\{\{1\}\}$  is in  $\mathcal{M}_1$ . For two sites, two possible partitions are available  $\{\{1, 2\}\}$  and  $\{\{1\}, \{2\}\}$ . For three sites, five par-

titions are available. The enumeration of  $\#\mathcal{M}_n$  is well understood and uses the Bell numbers (8). The sequence of the Bell numbers begins 1, 2, 5, 15, 52, 203, 877, 4140, 21147.

**Definition 7** The type of match of an  $n$ -tuple of observations  $\mathbf{y} = (y_1, \dots, y_n)$  is given by  $C(\mathbf{y}) \in \mathcal{M}_n$  where  $C(\mathbf{y}) = \mathbf{M} = \{M_1, \dots, M_m\}$  is the partition of the first  $n$  integers which satisfies  $(y_i = y_j) \Leftrightarrow i, j \in M_k$  for some  $k \in [1, 2, \dots, m]$ . That is,  $\mathbf{M}$  is the partition chosen such that any two site indices are in the same block within  $\mathbf{M}$  if and only if the observations in  $\mathbf{y}$  at those sites are equal.

It can clearly be seen that  $C(\mathbf{y})$  is uniquely specified by this definition. To continue with the earlier example, if

$\mathbf{y} = (\text{A123XYZ}, \text{A123XYZ}, \text{C789ABC})$  then

$$C(\mathbf{y}) = \{\{1, 2\}\{3\}\}$$

and if

$\mathbf{y} = (\text{A123XYZ}, \text{A123XYZ}, \text{A123XYZ})$  then

$$C(\mathbf{y}) = \{\{1, 2, 3\}\}.$$

It must now be shown that  $C(\mathbf{y})$  works as a representation of the type of match in a consistent way with the relationship  $\sim$  given by Definition 4.

**Theorem 8** For  $n$ -tuples of observations  $\mathbf{y} = (y_1, \dots, y_n)$  and  $\mathbf{z} = (z_1, \dots, z_n)$  then

$$C(\mathbf{y}) = C(\mathbf{z}) \text{ if and only if } \mathbf{y} \sim \mathbf{z}.$$

**PROOF.** Let  $\mathbf{M}_y = C(\mathbf{y})$  and  $\mathbf{M}_z = C(\mathbf{z})$ . First it must be shown that  $(\mathbf{y} \sim \mathbf{z}) \Rightarrow (\mathbf{M}_y = \mathbf{M}_z)$ . This follows trivially. Since  $(y_i = y_j) \Leftrightarrow (z_i = z_j)$  then if  $i, j$  are in the same set in  $\mathbf{M}_y$  they must be in the same set in  $\mathbf{M}_z$  and if they are in different sets in  $\mathbf{M}_y$  they must be in different sets in  $\mathbf{M}_z$ . As all integers from one to  $n$  appear once each in both partitions then it must be the case that  $\mathbf{M}_y = \mathbf{M}_z$ .

Similarly it must be shown that  $(\mathbf{M}_y = \mathbf{M}_z) \Rightarrow (\mathbf{y} \sim \mathbf{z})$ . A very similar argument applies. If  $i, j$  are in the same set in  $\mathbf{M}_y$  (and therefore in  $\mathbf{M}_z$ ) then  $y_i = y_j$  and also  $z_i = z_j$  if they are in different sets then  $y_i \neq y_j$  and also  $z_i \neq z_j$ . Therefore  $(y_i = y_j) \Leftrightarrow (z_i = z_j)$  and hence  $\mathbf{y} \sim \mathbf{z}$ .

It is useful at this point to define a shorthand notation for the type of match of most interest, that where the observations are the same at every site.

**Definition 9** Let  $\mathbf{A}_n \in \mathcal{M}_n$  represent a true match, that is the type of match where the same

observation is made over all  $n$  sites. Therefore,  $\mathbf{A}_n = \{\{1, 2, \dots, n\}\}$ .

## 5. Introducing false matching into the framework

So far the false match problem has been ignored and it has been assumed that for a given  $n$ -tuple of observations  $\mathbf{y} = (y_1, \dots, y_n)$  then the relation  $y_i = y_j$  can be taken at face value. However, the original problem was that, in licence plates, partial observations can lead to two distinct individuals being confused. In order to capture this in the described framework, a partial ordering will be introduced on the set  $\mathcal{M}_n$  and this will then be related to the partial observations. [It is somewhat unfortunate that this paper uses the phrase ‘‘partial plate survey’’ from transportation and the term ‘‘partial ordering’’ from set theory. These terms should not be confused.]

The next step is to introduce a partial ordering on the set  $\mathcal{M}_n$ . It will be seen in the next section how this relates to the false matching problem.

**Definition 10** For two partitions

$$\mathbf{M} = \{M_1, \dots, M_m\} \in \mathcal{M}_n$$

and

$$\mathbf{M}' = \{M'_1, \dots, M'_{m'}\} \in \mathcal{M}_n$$

a partial ordering  $\succsim$  is given by,

$$\mathbf{M} \succsim \mathbf{M}' \text{ if and only if } (i, j \in M_k) \Rightarrow (i, j \in M'_l),$$

for some  $k$  and  $l$ . Put more simply,  $\mathbf{M} \succsim \mathbf{M}'$  if whenever  $i$  and  $j$  are in the same set within  $\mathbf{M}$  then they are also in the same set within  $\mathbf{M}'$ .

The symbol  $\succ$  will be used to mean strictly succeeds. That is  $\mathbf{x} \succ \mathbf{y}$  means  $\mathbf{x} \succsim \mathbf{y}$  and  $\mathbf{x} \not\sim \mathbf{y}$ . The symbol  $\succ\triangleright$  will be used to mean immediate successor that is, if  $\mathbf{x} \succ\triangleright \mathbf{z}$  then  $\mathbf{x} \succ \mathbf{z}$  but there is no  $\mathbf{y}$  such that  $\mathbf{x} \succ \mathbf{y} \succ \mathbf{z}$ . The symbols  $\succ$ ,  $\succsim$  and  $\prec\prec$  will have their obvious meanings.

It can be trivially shown that this relation meets the formal requirements for a partial ordering. It should also be noted that this relation is extremely close to the original equivalence relation but with the implication going in one direction only. It can also be shown that under this partial ordering then  $\#\mathbf{M}$  the number of sets (blocks) in  $\mathbf{M} \in \mathcal{M}_n$  is a consistent enumeration of  $\mathcal{M}_n$ .

A Hasse diagram is a way of visualising a partially ordered set. A Hasse diagram is constructed by plotting a partially ordered set  $S$  graphically in such a way that for all  $\mathbf{x}, \mathbf{y} \in S$  if  $\mathbf{x} \prec \mathbf{y}$  then  $\mathbf{x}$  is further to the bottom of the diagram than  $\mathbf{y}$ . An arrow is drawn in a Hasse diagram from  $x$  to  $y$  if  $x \succ\triangleright y$ . Figure 1 shows the Hasse diagram of  $\mathcal{M}_4$  with the partial ordering given by the previous definition.

**Definition 11** Given an  $n$ -tuple of observations  $\mathbf{y} = (y_1, \dots, y_n)$ , let  $\mathbf{y}^* = (y_1^*, \dots, y_n^*)$  represent the partial observation formed from  $\mathbf{y}$ . Since a partial observation can cause distinct individuals to appear the same but cannot cause the same individual to appear distinct at different sites then the following relation holds,

$$(y_i = y_j) \Rightarrow (y_i^* = y_j^*).$$

This star notation will also be used to distinguish the set of all possible partial observations in the data  $\mathcal{S}^*$  and, in general, to distinguish functions which apply to partial data rather than the full data.

Note that this is the only assumption so far made about the nature of the partial observation. In licence plate surveys then the choosing of which part of a plate to survey needs to be made with reference to the particular format of plate to be observed. Consider the observations from the earlier example. If  $\mathbf{y} = (\mathbf{A123XYZ}, \mathbf{A123XDR}, \mathbf{C789ABC})$  then a standard way to make partial observations on this type of plate is to collect only the first letter and the digits. Therefore  $\mathbf{y}^* = (\mathbf{A123}, \mathbf{A123}, \mathbf{C789})$ . Note that  $C(\mathbf{y}) \neq C(\mathbf{y}^*)$  since  $y_1 \neq y_2$  but  $y_1^* = y_2^*$ . The way that  $C(\mathbf{y})$  can change when only a partial observation is made is given by the next theorem.

**Theorem 12** If  $\mathbf{y} = (y_1, \dots, y_n)$  is an  $n$ -tuple of observations then

$$C(\mathbf{y}^*) \preccurlyeq C(\mathbf{y}).$$

**PROOF.** Let  $\mathbf{M} = (M_1, \dots, M_m) = C(\mathbf{y})$  and  $\mathbf{M}' = (M'_1, \dots, M'_{m'}) = C(\mathbf{y}^*)$ . The theorem follows trivially from the relation given in 11. If  $i, j \in M_k$  for some  $k$  then  $y_i = y_j$  and hence  $y_i^* = y_j^*$  which in turn implies,  $i, j \in M'_l$  for some  $l$ . Therefore,  $(i, j \in M_k) \Rightarrow (i, j \in M'_l)$  which is the condition for the partial ordering.

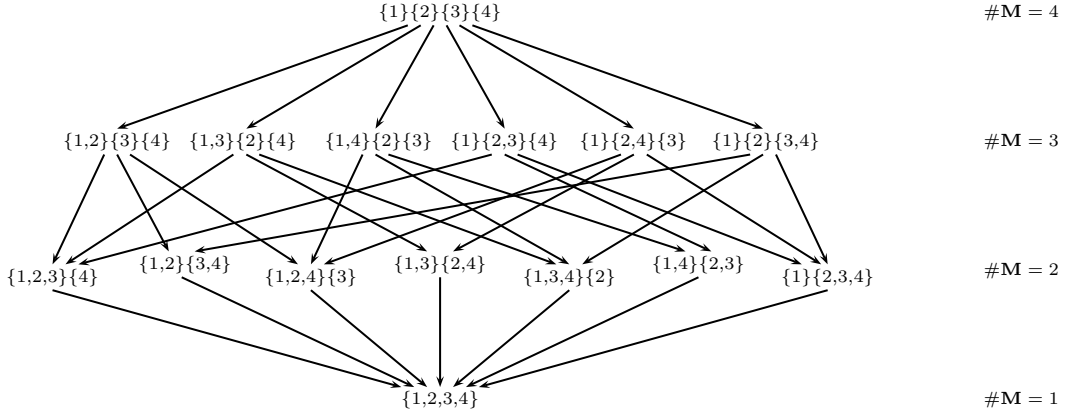


Fig. 1. Hasse diagram for  $\mathcal{M}_4$ .

From this theorem, it can be seen that when only partial data is available, the type of match of the partial observation may change only in a given way. Specifically, the type of match of the partial data can be the same as that of the full data or any type of match available by following down the arrows on the Hasse diagram.

Next, some counting functions are defined — these are used to enumerate the number of matches in the data which are different types of match.

**Definition 13** Let  $\mathbf{y}$  be an  $n$ -tuple of observations and  $\mathbf{M} \in \mathcal{M}_n$  be a type of match. The exact matching function for an observation  $\mathbf{y}$  is defined by,

$$x(\mathbf{y}, \mathbf{M}) = \begin{cases} 1 & \text{if and only if } C(\mathbf{y}) = \mathbf{M} \\ 0 & \text{otherwise .} \end{cases}$$

**Definition 14** Let  $\mathbf{M} \in \mathcal{M}_n$  be a type of match. The exact matching function for  $\mathcal{S}$  the set of all observations is given by,

$$X(\mathbf{M}) = \sum_{\mathbf{y} \in \mathcal{S}} x(\mathbf{y}, \mathbf{M}).$$

It can be readily seen that  $X(\mathbf{M})$  is the number of  $n$ -tuples  $\mathbf{y} \in \mathcal{S}$  which have a type of match  $C(\mathbf{y}) = \mathbf{M}$ . It can be further seen that the original problem of counting the number of individuals seen at all of  $n$  sites is the problem of evaluating  $X(\mathbf{A}_n)$ .

**Definition 15** Let  $\mathbf{y}$  be an  $n$ -tuple of observations and  $\mathbf{M} \in \mathcal{M}_n$  be a type of match. The relaxed matching function for an observation is defined by,

$$r(\mathbf{y}, \mathbf{M}) = \begin{cases} 1 & \text{if and only if } C(\mathbf{y}) \preceq \mathbf{M} \\ 0 & \text{otherwise .} \end{cases}$$

Equivalently,

$$r(\mathbf{y}, \mathbf{M}) = \sum_{\mathbf{M}' \preceq \mathbf{M}} x(\mathbf{y}, \mathbf{M}').$$

**Definition 16** Let  $\mathbf{M} \in \mathcal{M}_n$  be a type of match. The relaxed matching function for  $\mathcal{S}$  the set of all observations is given by

$$R(\mathbf{M}) = \sum_{\mathbf{y} \in \mathcal{S}} r(\mathbf{y}, \mathbf{M}).$$

Equivalently,

$$R(\mathbf{M}) = \sum_{\mathbf{M}' \preceq \mathbf{M}} X(\mathbf{M}').$$

It should be noted in passing that  $R(\mathbf{A}_n) = X(\mathbf{A}_n)$  since there are no  $\mathbf{M} \prec \mathbf{A}_n$ .

## 6. Solving the false match problem

In order to solve the false match problem, it is necessary to prove some simple lemmas which relate these counting functions. The main goal here is to estimate  $X(\mathbf{A}_n)$  (the number of  $n$ -tuples representing the same individual at all  $n$  sites) in terms of the partial data  $\mathcal{S}^*$ . The second goal is to do this in a way which does not involve investigating every single possible  $n$ -tuple. The reason for this is

that a realistic size for a traffic survey is of the order of one thousand vehicles. If there are six sites, then there are  $1000^6$  tuples to investigate and this would be far too slow computationally.

**Lemma 17** *Any exact matching function can be expressed in terms of relaxed matching functions and “lower” exact matching functions.*

$$X(\mathbf{M}) = R(\mathbf{M}) - \sum_{\mathbf{M}' \prec \mathbf{M}} X(\mathbf{M}').$$

**PROOF.** This follows trivially from Definition 16.

This expression can be used recursively so that any  $X(\mathbf{M})$  can be expressed as a function of  $R(\mathbf{M}')$  for all  $\mathbf{M}' \prec \mathbf{M}$ . The lemma can be thought of as being a version of the inclusion/exclusion principle for partitions of the integers under this partial ordering.

**Definition 18** *Let  $M = \{m_1, \dots, m_l\}$  be a set of integers, such that  $m_i \in \{1, 2, \dots, n\}$  for all  $i$ . Let  $\mathcal{S}'$  be the set of  $l$ -tuples of observations formed by the cartesian product,*

$$\mathcal{S}' = S_{m_1} \times S_{m_2} \times \dots \times S_{m_l}.$$

*In other words,  $\mathcal{S}'$  is the set of  $l$ -tuples of observations over some subset of the original sites. Then define,*

$$T(M) = X(\mathbf{A}_l),$$

*where the exact match  $X(\mathbf{A}_l)$  is in this case over the  $l$ -tuples in  $\mathcal{S}'$  rather than the  $n$ -tuples in  $\mathcal{S}$ . In other words,  $T(M)$  is the number of individuals seen at all sites in the set  $M$ .*

Note that, It can be easily seen that the problem of evaluating  $T(M)$  is either exactly the same as the original problem, if  $M = \{1, 2, \dots, n\}$  or it is a sub problem over a reduced number of sites. If  $M$  has a single member  $M = \{m\}$  then  $T(M)$  is simply the number of observations in set  $S_m$  that is,  $T(\{m\}) = \#S_m$ .

**Lemma 19** *The relaxed matching function  $R(\mathbf{M})$  where  $\mathbf{M} = \{M_1, \dots, M_m\} \in \mathcal{M}_n$  can be expressed as a product of exact matches over subsets of sites using the expression,*

$$R(\mathbf{M}) = \prod_{i=1}^m T(M_i).$$

**PROOF.** Clearly, for an  $n$ -tuple of observations  $\mathbf{y} = (y_1, \dots, y_n)$  then,

$$r(\mathbf{y}, \mathbf{M}) = \begin{cases} 1 & \text{if for all } i, j, k \text{ then} \\ & (i, j \in M_k) \Rightarrow (y_i = y_j) \\ 0 & \text{otherwise.} \end{cases}$$

Therefore,

$$\sum_{\mathbf{y} \in \mathcal{S}} r(\mathbf{y}, \mathbf{M}) = \#\{\mathbf{y} \in \mathcal{S} : (i, j \in M_k) \Rightarrow (y_i = y_j) \text{ for all } i, j, k\}.$$

The left hand side of this is simply  $R(\mathbf{M})$  as required. Since  $\mathcal{S}$  is the cartesian product then it can be seen that those  $\mathbf{y} \in \mathcal{S}$  which meet the condition are those which are picked out by  $T(M_i)$  and therefore the right hand side is  $\prod_{i=1}^m T(M_i)$  as required.

Note that if  $\mathbf{M} = \mathbf{A}_n$  then this expression simply says  $R(\mathbf{A}_n) = T(\{1, 2, \dots, n\}) = X(\mathbf{A}_n)$ . In all other cases, this allows a relaxed matching function to be expressed as a product of exact matching functions over a subset of the original sites.

**Definition 20** *The probability  $p(i)$  where  $i \in \{1, 2, \dots, n\}$  is the probability that, given that  $i$  observed individuals are all different in the full observation, they will all be the same in the partial observation. For the method described to work, this  $p(i)$  must be independent of the sites at which the vehicles are observed. By convention,  $p(1) = 1$ .*

It should be noted that this definition does place some restrictions on the type of data which can be analysed by this method and which types of partial observations are suitable. A discussion of  $p(i)$  in the context of licence plate observations follows this section. It is likely that other formulations of this problem would be possible if  $p(i)$  varies with the sites considered.

**Lemma 21** *An unbiased estimator  $\hat{t}$  for  $X(\mathbf{A}_n)$  is given by,*

$$\hat{t} = X^*(\mathbf{A}_n) - \sum_{\mathbf{M} \succ \mathbf{A}_n} p(\#\mathbf{M})X(\mathbf{M}).$$

**PROOF.** The quantity  $X^*(\mathbf{A}_n)$  is equal to  $X(\mathbf{A}_n)$  plus all those  $n$ -tuples of observations



which are false matches. Each element of the sum represents the number of false matches arising from a given type of match. Writing this out formally,

$$\hat{t} = X^*(\mathbf{A}_n) - \sum_{\mathbf{M} \succ \mathbf{A}_n} \mathbb{E}[\#\{\mathbf{y} \in \mathcal{S} : C(\mathbf{y}^*) = \mathbf{A}_n, C(\mathbf{y}) = \mathbf{M}\}].$$

The set  $\{\mathbf{y} \in \mathcal{S} : C(\mathbf{y}^*) = \mathbf{A}_n, C(\mathbf{y}) = \mathbf{M}\}$  is the set of  $n$ -tuples in the data  $\mathcal{S}$  which are a match of type  $\mathbf{M}$  in the complete data but appear to be a match of type  $\mathbf{A}_n$  in the partial data  $\mathcal{S}^*$ . Now, the number of distinct individuals in this  $n$ -tuple must be  $\#\mathbf{M}$ . Therefore,

$$\mathbb{P}[C(\mathbf{y}^*) = \mathbf{A}_n | C(\mathbf{y}) = \mathbf{M}] = p(\#\mathbf{M}).$$

Bayes theorem gives,

$$\begin{aligned} \mathbb{P}[C(\mathbf{y}^*) = \mathbf{A}_n, C(\mathbf{y}) = \mathbf{M}] &= p(\#\mathbf{M})\mathbb{P}[C(\mathbf{y}) = \mathbf{M}] \\ &= \frac{p(\#\mathbf{M})X(\mathbf{M})}{\#\mathcal{S}}. \end{aligned}$$

Hence, the expected number of false matches arising from each type of match can be given by,

$$\begin{aligned} \mathbb{E}[\#\{\mathbf{y} \in \mathcal{S} : C(\mathbf{y}^*) = \mathbf{A}_n, C(\mathbf{y}) = \mathbf{M}\}] &= \#\mathcal{S}\mathbb{P}[C(\mathbf{y}^*) = \mathbf{A}_n, C(\mathbf{y}) = \mathbf{M}] \\ &= p(\#\mathbf{M})X(\mathbf{M}), \end{aligned}$$

and the lemma follows immediately.

It may not be immediately obvious that Lemmas 21, 17 and 19 together allow an unbiased estimate of the number of true matches, from the partial plate data (assuming that the  $p(i)$  are known). First, looking at Lemma 21, the quantity  $X^*(\mathbf{A}_n)$  can be simply enumerated by computer in the partial data. Therefore, this lemma allows an unbiased estimate of the number of matches in the complete data if an unbiased estimate of  $X(\mathbf{M})$  can be found for all  $\mathbf{M} \succ \mathbf{A}_n$ . Now, Lemma 17 allows  $X(\mathbf{M})$  to be expressed as a sum of  $R(\mathbf{M}')$  for all  $\mathbf{M}' \lesssim \mathbf{M}$ . Lemma 19 allows those  $R(\mathbf{M}')$  to be either equal to the original required quantity  $X(\mathbf{A}_n)$  or to be expressed in terms of a product involving subproblems on a reduced number of sites. Hence, computer algebra can be used to give an equation which is in terms of  $X(\mathbf{A}_n)$  (the quantity desired),  $X^*(\mathbf{A}_n)$  (measurable on the data),  $p(i)$  (assumed to be known) and  $T(M)$  (which is a subproblem

of the original problem with a reduced number of sites). The computer can then be used to recursively solve the subproblem which has already been shown to be trivial for just one site. An expanded description of this solution process is given in (1, Chapter 4).

### 6.1. Estimating the probability of false matches

The method described here relies on a good estimate of  $p(i)$  and also on the assumption that this does not vary by the sites chosen. The specific details of British licence plates are not of general interest (and it should again be stressed that the method discussed here is general and not limited just to specific types of licence plate survey, indeed it could be used for any type of data collection where the restrictions on  $p(i)$  are met). However, illustrating how  $p(i)$  can be estimated in a practical case might be of interest and illuminate how the method was applied in real life. More details on this can be found in (1, Chapter 5).

Two methods of estimating  $p(i)$  are practical. If the distribution of the vehicle types can be calculated then an analytical approach is possible. Let there be  $N$  vehicle types which are distinguishable in the partial observations and let  $f_j$  be the proportion of the vehicle fleet which is of type  $j$  (assume the membership of each type is relatively large). Therefore,  $p(i)$  is approximately given by,

$$p(i) = \sum_{j=1}^N f_j^i.$$

In the case of the old style British licence plates discussed, the distribution of the digits is almost a flat distribution from 1 to 999. The distribution of the year letters is more complex and can be estimated from consideration of the data. Therefore,  $f_j$  can be calculated for each possible partial observation and hence  $p(i)$ .

An alternative method is to estimate  $p(2)$  by finding two sites which are so far separated geographically that no vehicle could be seen at both. Any vehicle seen at both must be a false match and therefore if there are  $x$  observed matches in the partial data and then  $p(2) = x/(\#\mathcal{S}_1\#\mathcal{S}_2)$ . Similarly  $p(3)$  can be estimated by finding three such geo-

graphically distant sites. Higher order  $p(i)$  can be estimated with reference to the previous method or by assuming a functional form for the fall off. An estimate of  $p(2) = 7.4 \times 10^{-6}$  was given in (1, Chapter 5) for licence plate data of the type discussed.

## 7. Results on simulated data

Table 7 shows simulation results for between two and six observation sites. These could be thought of as one site observed on several days, or six sites observed on several different days. Num. Veh. refers to the total number of observations at each of the sites (in these simulations, there are the same number of vehicles in each data set). The five columns of the form  $1 - n$  refer to the number of vehicles which genuinely went from site one to site  $n$  visiting all sites in between. If this column is blank it means that there was no site  $n$ . For example, if  $1 - 2 = 100$ ,  $1 - 3 = 200$  and  $1 - 4$  is blank. This means that 100 vehicles travelled between site one and site two, 200 vehicles travelled between sites one, two and three and there were only three sites. Note that these are cumulative so that if  $1 - 2 = 20$  and  $1 - 3 = 10$  this means that 30 vehicles in total went from site one to site two and ten of them continued to site three. Thus the first experiment is two sites, 1000 vehicles at each for which there were ten vehicles which were genuinely seen at both sites. In every experiment, the number of different vehicle types was set at 10,000 with a flat distribution (equal numbers of vehicles seen at each site). Note that the simplifying assumptions of a flat distribution and the same number of vehicles at each site are simply there to make the experiment easier to understand rather than being necessary for the method to work. It should be clear that the desired answer from the correction process is the rightmost figure in these columns.

Each experiment is repeated twenty times with simulated data being generated anew each time. The correction process has no random element and will always give the same result for the same data. The mean raw number of matches is given — this is the total number of n-tuples which were seen to

have the same value for each observation at every site (averaged over the twenty simulation runs). Because of the combinatorial nature of the procedure, this could, in principle, be much larger than the number of vehicles in any of the data sets (since it counts any n-tuple). The sample standard deviation ( $\sigma$ ) is given for the raw matches. The mean estimated correct number of matches is then given (again averaged over the twenty simulations). The sample standard deviation  $\sigma$  is then given for the twenty corrected matches. It is clear that the most important test is that the mean corrected number of matches is as near to correct as possible. However, it should also be kept in mind that in reality, a researcher could only run the matching procedure once on any given set of data so it is also important that  $\sigma$  is as low as possible. A significant improvement to the method would be to estimate the variance as well as producing the mean in order that the researcher could have some idea as to the likely accuracy of the corrected results.

The first five rows are all results on just two test sites. This procedure is not the ideal one to use for estimates on matches between just two sites and the work of other authors in the field should be used in such a circumstance. However, these results are included here for completeness. In the first experiment, the average number of raw matches over the twenty runs is 111.4. The average number of corrected matches is 100 less than this (11.4). This is close to the correct answer of 10. However, it should be noticed that the  $\sigma$  is high in comparison to the actual answer. In this case, the  $\sigma$  is 8.5 which is of the same order of magnitude as the answer. This is to be expected since we are looking for only 10 true matches in over 110 observed matches. If we increase the number of vehicles to 2000 then, as would be expected, the number of false matches goes up (to approximately 400) and the  $\sigma$  also rises (to almost 20).

The next five rows of results are all over three sites. In the first of these, 10 vehicles travel between all three and all other matches are coincidence. 1000 vehicles are observed at all sites. The mean corrected match across all sites 9.3 is close to the actual answer of 10 and the  $\sigma$  is lower than in the two site case. However, when the same experiment is run with 500 vehicles travelling from sites one to

two in addition to 10 vehicles travelling from sites two to three, the  $\sigma$  increases markedly (it almost doubles). In all cases with three sites, the mean is a good estimate and the  $\sigma$  is generally low enough that a good estimate can be expected.

The next four rows of results are for experiments made over four sites. The first experiment has 100 vehicles which visit all four. The mean corrected match is 104 (very close) and the  $\sigma$  is only 22. It is hard to explain why this  $\sigma$  actually falls in the next experiment when more vehicles are genuinely seen in common between the other sites. In all cases the mean of the predictions is approximately correct (the worst performance being in the case of the fourth experiment when the mean was 106.1 not 100).

The next six rows of results are experiments made over five sites. Again, the mean corrected results are approximately correct. However, in the worst case, the mean is 11 too high and the  $\sigma$  in the results is 46.7 which is comparable to the level of the effect being observed. In this case approximately 120 false matches are being removed each time. However, previous experiments have been able to correct for a greater proportion of false matches with less  $\sigma$  in the result.

The final four rows of results are experiments over six sites. This was the largest number of sites for which it was practical to do runs of twenty or more simulations with the computer power available. Again, the mean corrected estimate of matches was nearly correct in all cases. The worst performance was an estimate of 92.2 (correct result 100). The  $\sigma$  was, however, relatively high. This was a surprise in some cases — particularly the first row of results where the mean number of false matches was only 21.2. In many senses, the worst results was the final one where a  $\sigma$  of 55.0 was given on an corrected prediction of only 101.3.

The time taken to do one run over six sites with one thousand pieces of data on each site was thirty seconds on a Celeron 366 computer running Debian Linux. Six sites with one thousand vehicles at each is a reasonable size for a typical traffic survey. It is practical (if time consuming) to do experiments on seven sites, even using such comparatively obsolete equipment. However, eight sites or more is probably too computationally expensive for the moment

and this is a limitation of the method outlined. The exact rate at which the computational requirements increase with the number of sites is hard to determine. It will relate to the Bell numbers, to the number of observations at each site and to the number of pairs of observations at each site pair.

The results given here are certainly consistent with the idea that the method gives an unbiased estimator for the true number of matches. In some experiments, there were problems with the standard deviation being higher than would be desirable in real cases. It is important to bear in mind that these were relatively extreme tests of the method since  $p(2)$  and  $p(3)$  were relatively low and the number of samples given were quite high. Often the method was attempting to predict only ten true matches in a number of observed matches which might be several hundred.

To test the method more fully, four very extreme tests were given. Each of these tests involved six sites at each of which one thousand vehicles were observed. Interacting flows were chosen to cause a large number of false matches in a diversity of ways. Because these experiments were chosen to cause a large number of false matches then one thousand runs of each experiment were performed. The averaged results are shown in Table 2.

In experiment one, five hundred vehicles travelled from one to five and five hundred from two to six. The remaining five hundred vehicles at sites one and six were appeared nowhere else. No vehicles made the complete journey. As can be seen, on average over seven hundred false matches were seen and the standard deviation between runs was extremely large. However, the mean was within twelve of the correct answer (zero) although the standard deviation was large. In such extreme circumstances, a single experiment would be next to useless but it is good evidence that the method was unbiased.

In experiment two, five hundred vehicles travelled from one to three. Five hundred vehicles travelled from four to six. Five hundred vehicles visited only odd numbered sites and five hundred vehicles visited only even numbered sites. In this experiment the corrected mean result was almost exact (within one) and the standard deviation was much lower than the other three experiments.

No. Veh.	1 - 2 1 - 3 1 - 4 1 - 5 1 - 6	Av. Raw Matches	$\sigma$ Raw Matches	Av. Cor. Matches	$\sigma$ Cor. Matches
1000	10	111.4	8.5	11.4	8.5
2000	10	411.8	19.5	11.8	19.5
1000	100	199.2	12.0	99.2	12.0
1000	200	302.3	7.7	202.3	7.7
1000	500	596.6	12.3	496.7	12.3
1000	0 10	21.9	4.6	9.3	3.3
1000	500 10	73.8	7.5	10.2	6.2
1000	100 100	152.1	8.5	101.9	7.5
1000	500 250	388.3	22.7	253.2	20.1
1000	0 500	667.2	24.9	506.0	22.3
1000	0 0 100	154.6	26.6	104.0	22.6
1000	100 100 100	164.4	11.4	97.7	9.3
500	100 100 100	140.7	19.3	105.8	17.4
1000	500 250 100	207.8	29.7	106.1	23.7
500	10 10 10 10	14.2	2.2	10.5	1.8
1000	10 10 10 10	17.4	4.1	9.4	2.8
500	50 50 50 50	71.3	14.3	47.8	12.3
500	100 100 100 100	151.9	26.9	92.0	22.3
1000	0 0 0 100	177.6	29.9	103.4	22.6
1000	100 100 100 100	222.2	61.5	111.0	46.7
1000	0 0 0 0 10	21.2	13.4	12.3	9.9
500	0 0 0 0 100	152.6	45.5	92.2	37.3
1000	0 0 0 0 100	214.6	58.0	103.5	40.2
1000	100 100 100 100 100	289.8	88.4	101.3	55.0

Table 1  
Simulation results — all performed over twenty runs with 10,000 distinct vehicle types.

Experiment Number	Expected Answer	Av. Raw Matches	$\sigma$ Raw Matches	Av. Cor. Matches	$\sigma$ Cor. Matches
1	0	739	305	11.9	196
2	0	110	45.5	-0.950	27.1
3	250	836	287	249	205
4	500	1920	531	496	356

Table 2  
Simulation results — all performed over one thousand runs with 10,000 distinct vehicle types.

In experiment three, two hundred and fifty vehicles travelled to all sites. Five hundred vehicles went from site one to three and five hundred from four to six. The remaining two hundred and fifty vehicles at each site visited only that single site. As can be seen, the corrected result is almost exactly correct although, again, the standard deviation is so high that a single reading would be worthless.

In experiment four, five hundred vehicles visited every site. Two hundred and fifty vehicles went from sites one to three. Two hundred and fifty vehicles went from sites four to six. Two hundred and fifty vehicles visited only sites one and two, two hundred and fifty vehicles visited only sites three and four and two hundred and two hundred and fifty vehicles visited only sites five and six. Again, the mean of all results is very close (within four vehicles) but the standard deviation is the highest yet seen. This is not surprising. The mean number of raw tuples of matches averaged nearly 2000, twice the number of vehicles at each site.

These four tests provide a convincing demonstration that the method is, indeed, unbiased as was shown by theory.

## 8. Conclusions

This paper presented a framework for analysis of surveys where matches are required over more than two data collection points. The framework given formalises the concept of a type of match using the concept of the equivalence class. Further a method is given for evaluating  $\mathcal{M}_n$  the set of all possible types of match over multiple data sets.

The framework given is then applied to the problem of false matches — which is put into the language of set theory using the concept of a partial ordering. It is shown how this partial ordering can be used to visualise, by means of a Hasse diagram, the ways in which false matches can occur in data observed at multiple sites. The framework was then used to design and implement an algorithm which was used to estimate the number of true matches in simulated data. The algorithm has also been tested on real data from partial plate surveys.

This algorithm was implemented and tested on

simulated data. The results show that the estimator seems to be unbiased and in the majority of cases tested the standard deviation on the results is low. The method is suitable for analysis of matches on data between three and seven test sites but becomes too computationally intensive after this point. A significant improvement to the method would be the estimation of a variance as well as a corrected number of matches. A potential weakness of the method is that it relies on good estimates for  $p(i)$ .

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